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Tofu Interconnect 2: System-on-Chip Integration of High-Performance Interconnect

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Fujitsu Limited

Introduction

- **Tofu interconnect** (Tofu1) was developed for the K computer
  - '**To**rus **fu**sion' derives from its network topology
- Introducing **Tofu interconnect 2** (Tofu2)
  - Designed for Fujitsu's next generation machine Post-FX10
  - SoC integration, improved link speed and new efficient functions

K computer

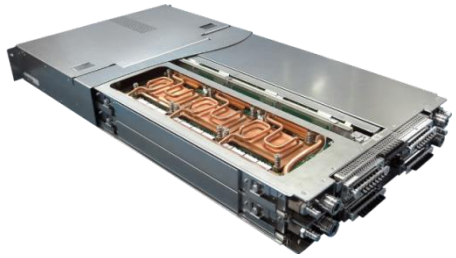
8 core

FX10

16 core

DDR3 SDRAM
Tofu interconnect

Post-FX10



32 core

Hybrid Memory Cube
Tofu interconnect 2

2010 **2012** **2015**

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- Network topology
- SoC integration
- Improved link speed
- New efficient functions
- Summary

Summary of Tofu Interconnect 2

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	Tofu1	Tofu2
Network topology	6D-Mesh/Torus	←
# of network interfaces	4	←
# of network links	10	←
Implementation	Discrete chip	SoC integration
Link speed	40 Gbps	100 Gbps
# of optical links	0	6 - 7
RDMA Put	✓	✓
RDMA Get	✓	✓
RDMA Atomic RMW		✓
Barrier synchronization	✓	✓
Non-blocking collective		✓
Memory bypass (sender)	✓	✓
Memory bypass (receiver)		✓

6D-Mesh/Torus network topology

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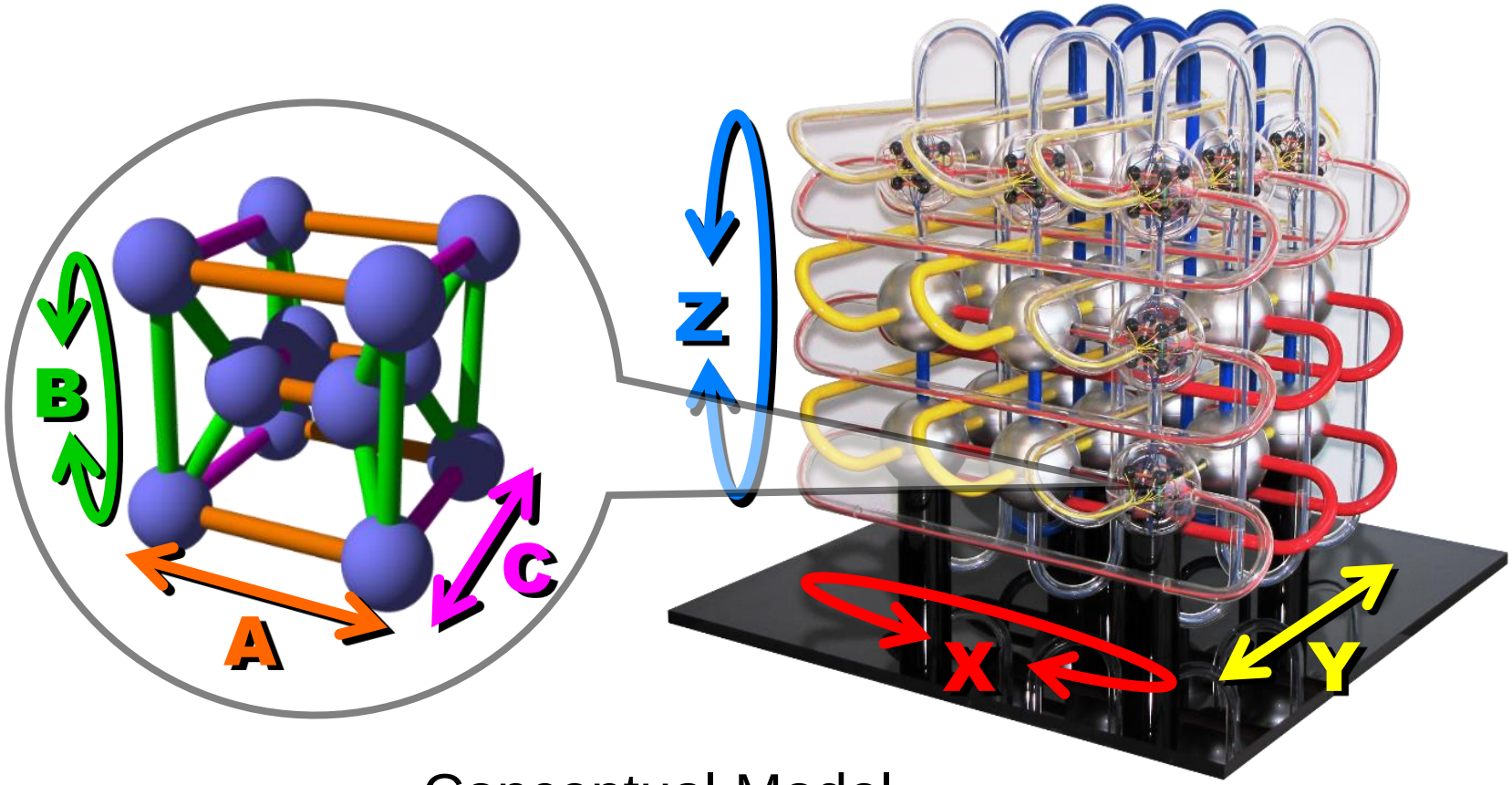
- Every 3-D Cartesian grid point embeds a 3-D structure
- Higher bisection bandwidth than 3D-Torus
- Virtual torus for topology-aware communication algorithms



Conceptual Model

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System-on-Chip Integration

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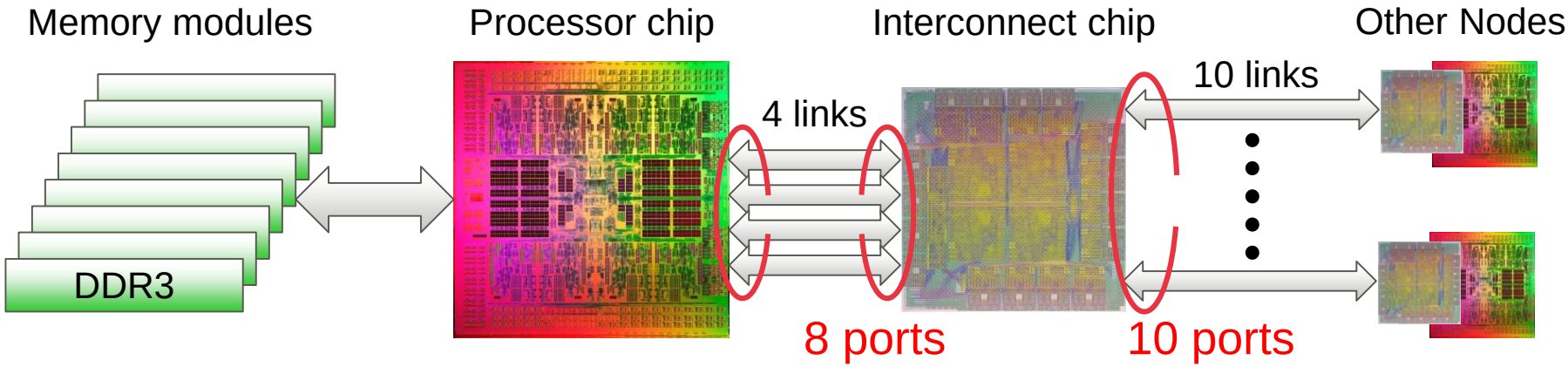
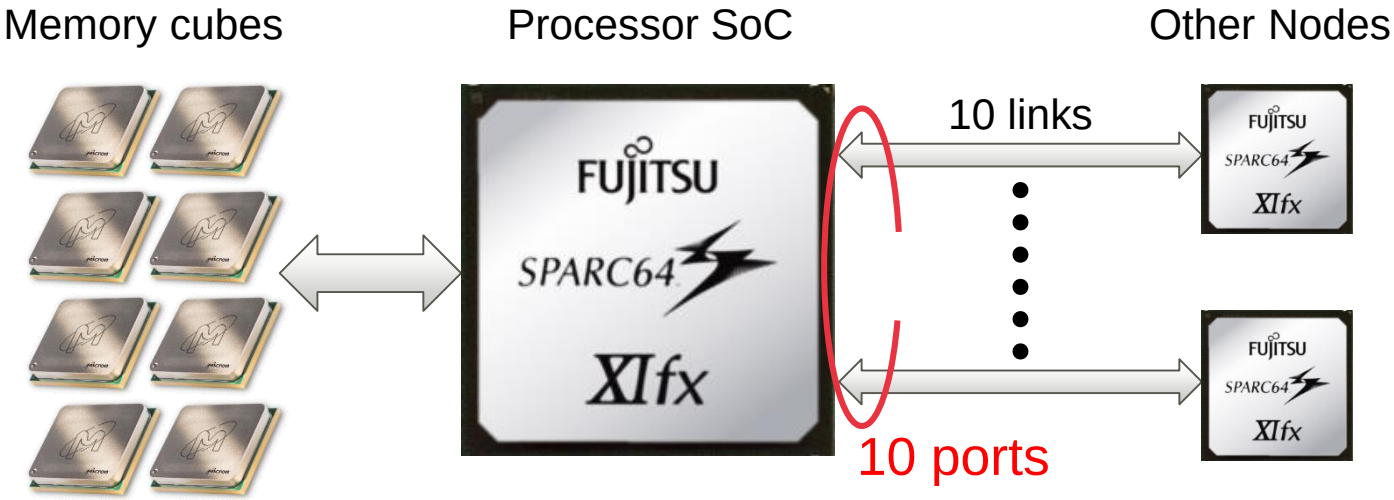
- Tofu1 was implemented as a discrete chip



- Tofu2 is integrated into a processor SoC



- Number of ports per node decreased from 18 to 10

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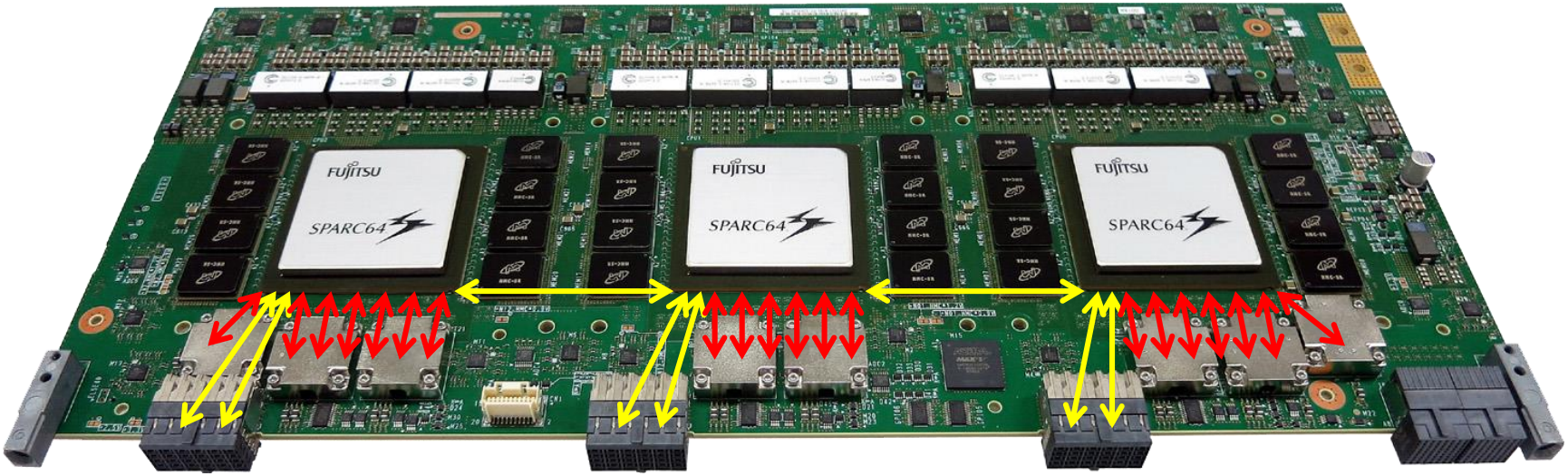
Transmission Technology

- Number of signals per link decreases from 8 to 4 lanes
  - Tofu1: 8 lanes × 18 ports = 144 lanes
  - Tofu2: 4 lanes × 10 ports = 40 lanes (limited pin-count)
- Link speed increases from 40 Gbps to 100 Gbps
  - By increasing data transfer rate fourfold from 6.25 to 25.78125 Gbps
- 2/3 of the links are optical
  - 1 out of 3 nodes uses **6 optical links** + **4 electrical links**
  - 2 out of 3 nodes use **7 optical links** + **3 electrical links**



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RDMA Atomic Read-Modify-Write

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- Atomically read, modify and write back remote data
  - Typical operations: compare-and-swap and fetch-and-add
  - Usage: software-based synchronization and lock-free algorithms
- Atomicity
  - Guaranteed by extending the coherency protocol of processor
    - not by each network interface
  - Strong atomicity: Any memory accesses cannot break atomicity
  - Mutual atomicity: Atomic operations of processor and Tofu2 mutually guarantee their atomicity
- The mutual atomicity enables an efficient implementation of unified multi-process and multi-thread runtime

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Cache Injection

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- Injecting received data into L2 cache directly
  - Bypassing main memory
- Injection flag On/Off is indicated by the sender
- Reduction of communication latency
  - The evaluations used the Verilog RTL codes for the production
  - Communication pattern: Ping-Pong of Put transfer

Injection flag	Estimated half round-trip latency
Off	0.87 usec
On	0.71 usec

0.16 usec reduction

- Harmless injection
  - Cache injection is only performed when cache hits and the line is in exclusive state
  - A cache line in exclusive state is highly likely to be polled by a corresponding processor core

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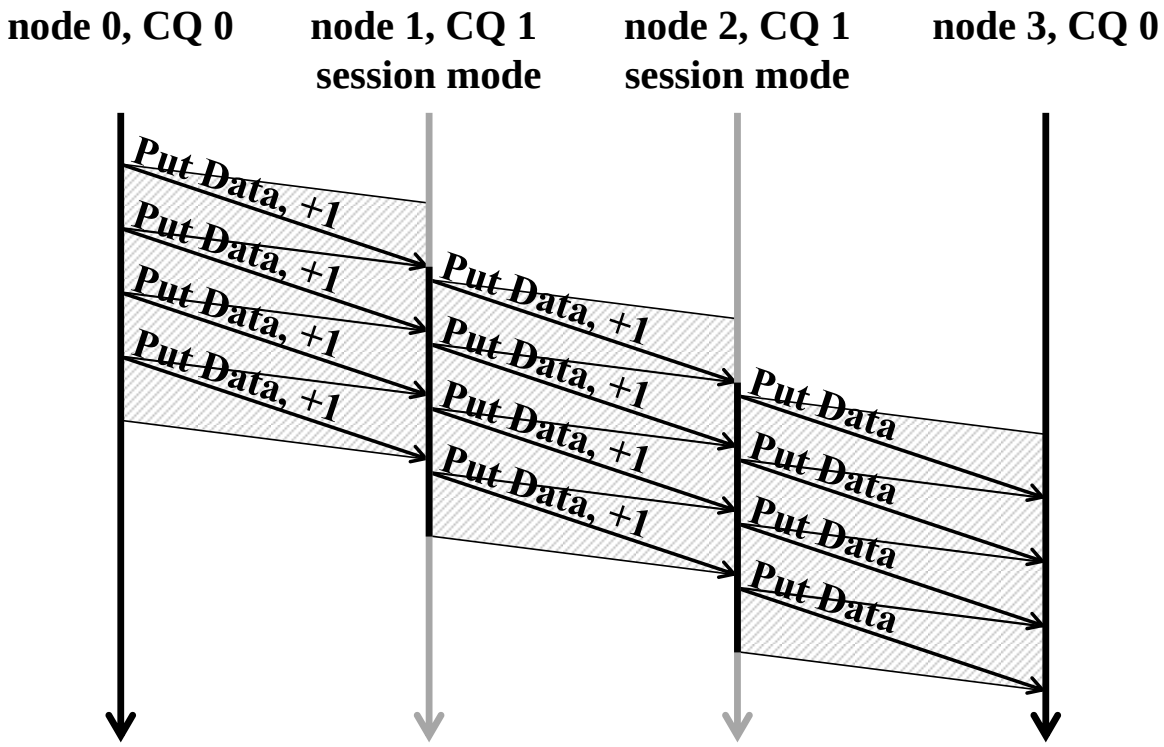
Session-mode Control Queue (CQ)

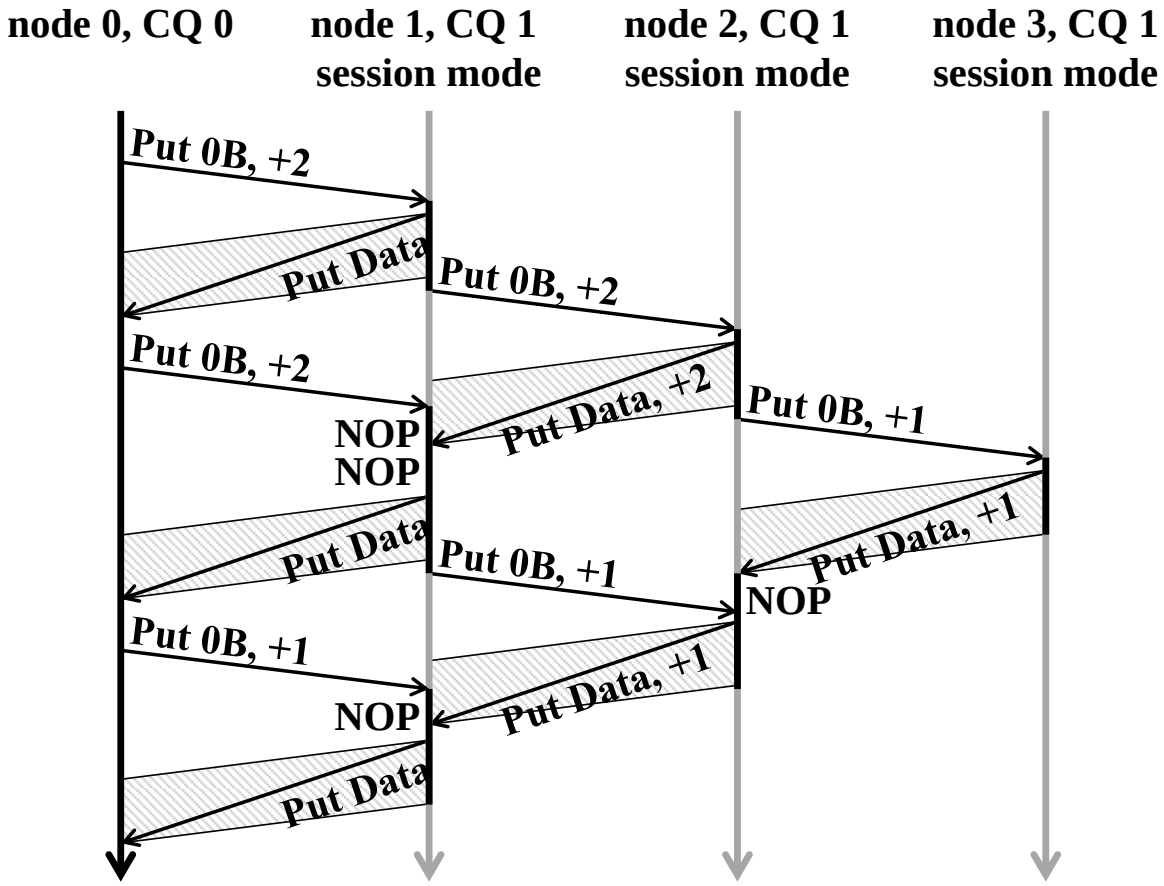
- Offloading a collective communication of long messages
- Command execution in a session-mode CQ is only advanced on a successful reception of Put transfer



Example of offloading pipelined Broadcast communication

Flexibility of Session-mode CQ

- Control flow can be branched or joined
  - Branch by advancing multiple commands
  - Join by enqueuing no operation commands



Example of offloading handshaking Gather communication

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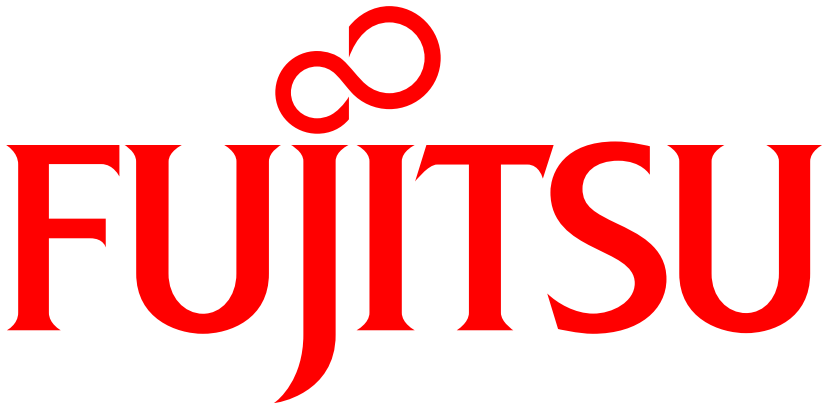
  - Designed for the next generation Post-FX10 machine
- System-on-chip integration
  - The number of link ports per node decreases from 18 to 10
- Link speed increases from 40 Gbps to 100 Gbps
  - 2/3 of the links are optical
- New efficient functions
  - Tofu2's atomic RMW operations guarantee the mutual atomicity that enables to unify multi-thread and multi-process shared variables
  - The cache Injection function bypasses main memory on a receiver side and reduces communication latency by 0.16 usec without cache pollution
  - The flexibility of session-mode CQ enables offloading various non-blocking collective communication algorithms



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