



THE POSSIBILITIES ARE INFINITE

XCAST6

--eXplicit Multicast on IPv6

----Jack Lü

Telecommunications Group

FUJITSU China LTD

<http://telecom.fujitsu.com>

2003.4





Outline



Motivation

Basic idea

Deployment method

- Semi-permeable capsule

Implementations

Trials

IETF Standardization



When can we truly start to use inter-domain multicast?

- For a decade and more, many Mbone talented researchers have worked in and out of IETF.
 - In the lab – we can make one easily.
 - On academic testbed nets – very hard but possible
 - In commercial Inter-domain net - never
- We want inter-domain multicast env. as easy as unicast we daily use, however...



Our approach


Focused on what we really expect for multicast

- Don't solve the whole multicast problems.

At first, think about what truly we need.

- Want to be broadcaster? – No.
- Want to deliver the Hollywood movies? - No.
- Want to communicate with my friends? – Yes!

Category of Multicast Applications

Focus! 

<p>Broadcast-like (one-to-many)</p> <ul style="list-style-type: none">• Multicast of IETF meetings• Broadcast of TV programs	<p>Narrowcast-like (a few-to-a few)</p> <ul style="list-style-type: none">• IP Telephony with conferencing• Video conferencing• Real-time collaborative applications• Multiparty networked games
---	---



Existing Multicast

(scales with number of receivers)



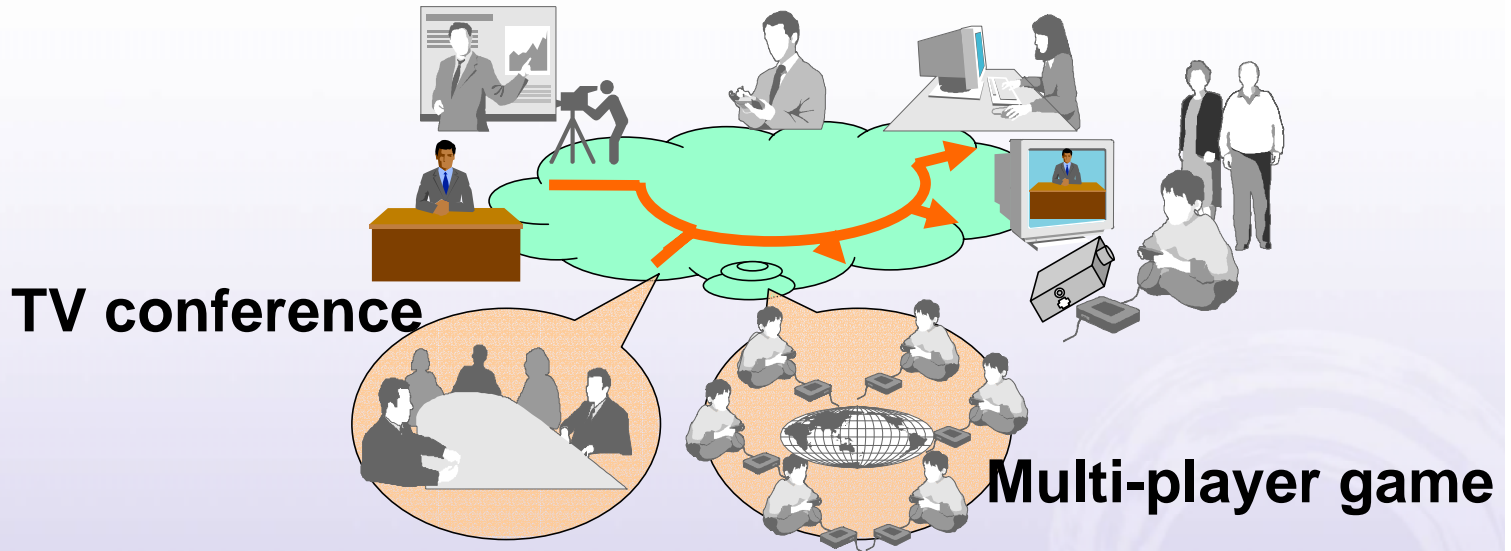
Small Group Multicast

(scales with number of sessions)

(Source Dirk Ooms in Alcatel.com)

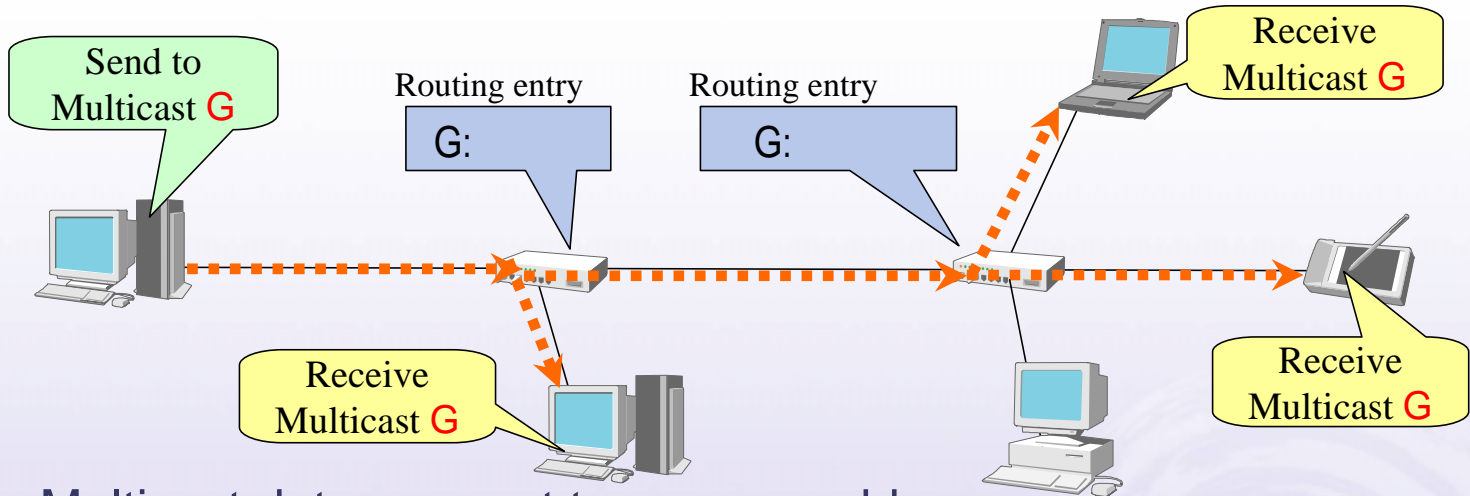


Goal: Narrowcast like multicast



To deliver for limited small number of nodes
Network must support very huge number of small groups.
Anybody can transmit from anywhere on the Internet

Existing Multicast mechanisms



Multicast data are sent to a group address.

All routers along the delivery path must maintain the status for each group.

Intermediate routers need to know where the sender is in order for new nodes to join a multicast group.

Receivers periodically send keep alive messages.



Scalability Problem of Existing Multicast

Protocol	Table Size/Cntl Msg
DVMRP	$O(G*S)$
PIM-SM	
- shared	$O(g)$
- short cut	$O(g*S)$
logical lower bound	$O(g)$

G: # of active groups
in the DVMRP domain.

g: # of groups
running on a router

S: # of source

Sola&Ohta “Scalability of Internet multicast Protocols”, Inet 98

For 1 million multicast groups, we must hold 1 million routing entries and process 1 million join/prune messages per min.



Outline

Motivation



Basic idea

Deployment method

- Semi-permeable capsule

Implementations

Trials

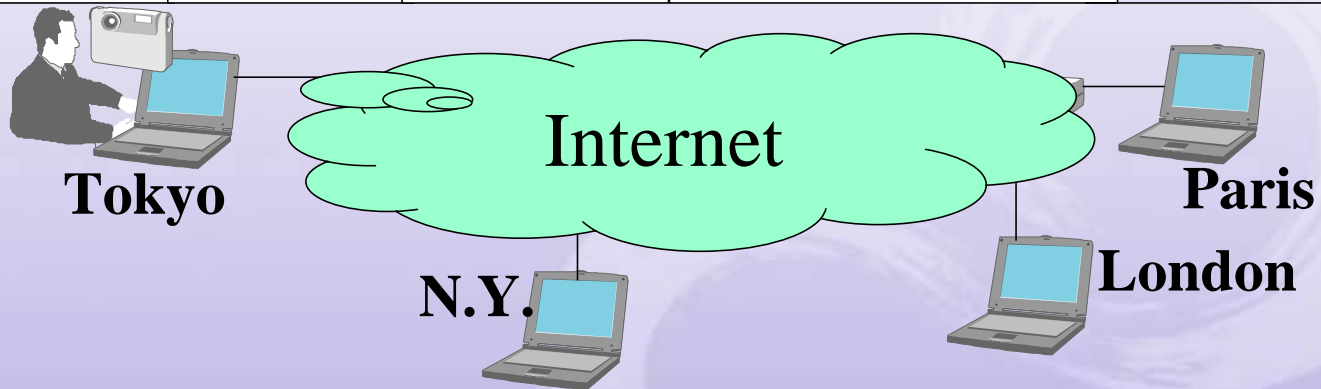
IETF Standardization

Main idea of XCAST6

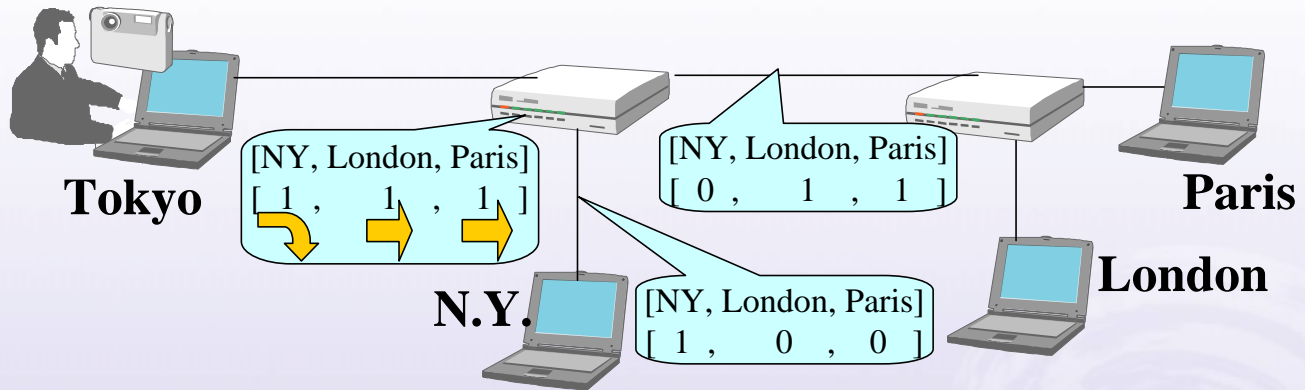


Instead of a group address, an explicit list of unicast destination addresses is stored in an optional IPv6 routing header.

IPv6 header SRC=Tokyo DST=N.Y.	Hop-byHop header TAIL=Paris	IPv6 header SRC=Tokyo DST=XCAST.	ROUTING header [N.Y., London, Paris] [1 , 1 , 0]	UDP header
--------------------------------------	-----------------------------------	--	--	---------------



Routing procedure



Look-up the next-hop for each address using a unicast routing table.

Bundle up destinations which have same next-hop

Advantages

XCAST6 can be delivered using only unicast routing information.

No need for

- a special multicast routing protocol
- maintaining multicast status on intermediate routers
- group address allocation
- sender location advertisement

Unlimited Scalability with respect to the number of groups

Advantages(Cont'd)

Explicit end-to-end control of multicast group membership.

- Senders can start transmission anytime without any signaling.
 - With existing multicast, receiver must join before transmission.
- Sender can change the group membership (destinations) per packet basis.
 - With existing multicast schemes, membership change is done by join/prune process, a complex process.

Disadvantages

Limit with number of receivers

- Logically, up to 126 destination in a IPv6 routing headers (8*256 octet).
- Actually, up to 15 destination for 1024 octets RTP video payload.
 - Ethernet MTU(1500 octet)
 - headers (XCAST, UDP, RTP)

Outline

Motivation

Basic idea

 Deployment method

- Semi-permeable capsule

Implementations

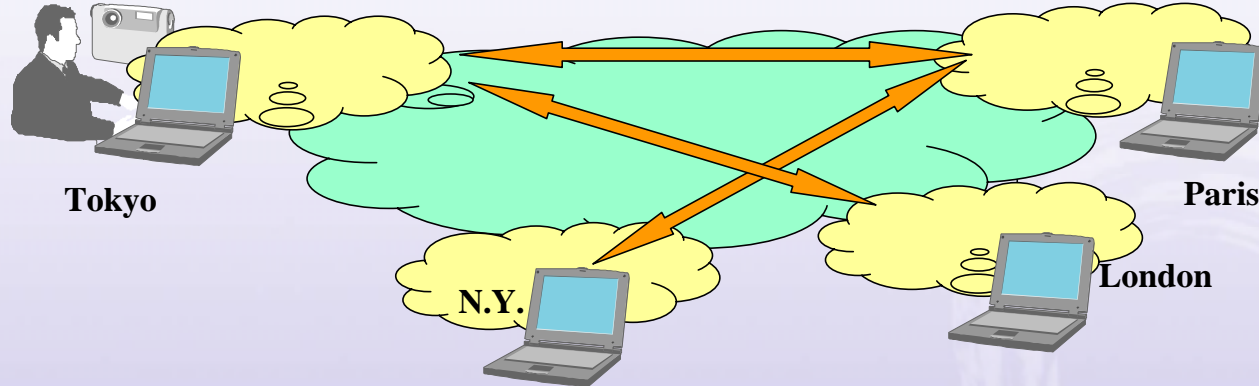
Trials

IETF Standardization



Deployment of existing multicast

Connecting multicast islands by IDMR and tunneling.



- negotiation between network operators
- complicated management
- encapsulation/peeling cost

**Big obstacle
for deployment**

Semi-permeable capsule

- The intermediate router which does not support XCAST6 treats a XCAST6 datagram as a regular unicast datagram.

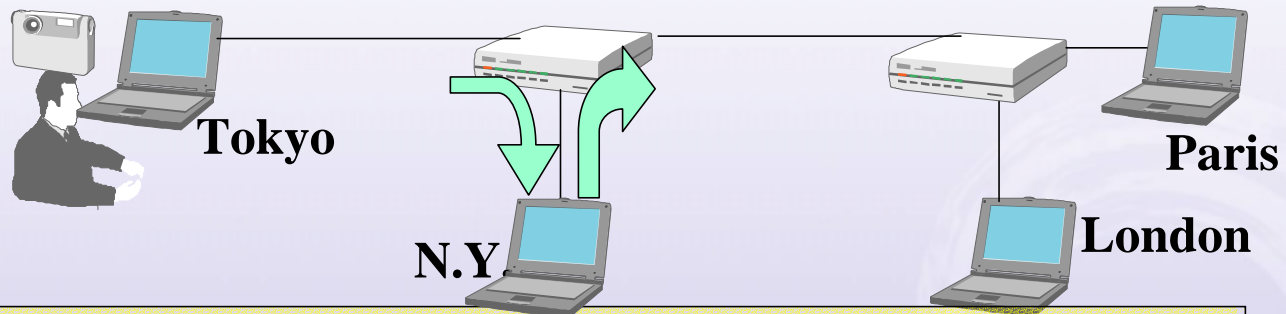
IPv6 header SRC=Tokyo DST=N.Y.	Hop-byHop header TAIL=Paris	IPv6 header SRC=Tokyo DST=XCAST.	ROUTING header [N.Y., London, Paris] [1 , 1 , 0]	UDP header
--------------------------------------	-----------------------------------	--	---	---------------

Temporal destination

Type prefix has '01' that means "ignore this option and forward" if router doesn't know this option.

Semi-permeable capsule(cont'd)

Even if non-XCAST6 routers are on the way, XCAST6 datagrams pass them once and turn back to next destination at next XCAST6 node.



- i. End node can transmit XCAST6 in any environment.
- ii. Installing more XCAST6 routers, path become optimized gradually.

Motivation

Basic idea

Deployment method

- Semi-permeable capsule

 **Implementations**

Trials

IETF Standardization

Implementations

WIDE project/FUJITSU Lab.

- OS: NetBSD 1.6, FreeBSD 4.6.2
- VIC (Video Conference) & RAT (Robust Audio Tool)
- <http://www.sourceforge.net/projects/xcast6>

ETRI/Soongsil University

- OS: Linux 2.4.18
- VIC & RAT
- <http://www.ipv6.or.kr/xcast/>

Outline

Motivation

Basic idea

Deployment method

- Semi-permeable capsule

Implementations

 Trials

IETF Standardization



Various meeting and events

Weekly WIDE XCAST WG meeting

- Discuss and steer this R&D activity itself.

Monthly BUGs(*BSD Users Groups) meeting

- For promotion into the open source community.

JP-BUGs (BSD Users Groups) meeting



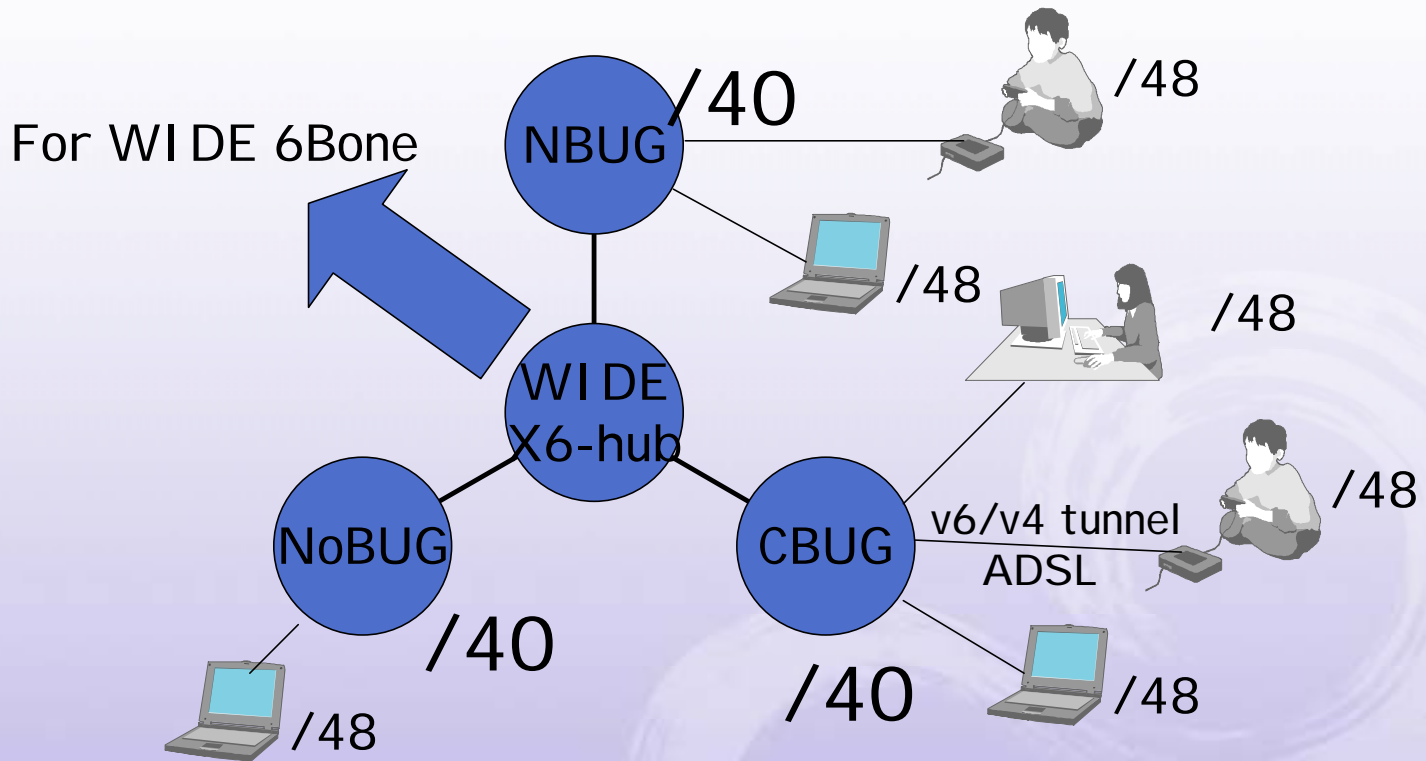
WIDE X6-Bone project

For more smooth XCAST6 delivery, more XCAST router must be deployed.

Make pseudo XCAST6 network using special pTLA space (3ffe:051b::/32).

Distribute **/40** for collaborators (BUGs, LUGs) and connect them for HUB in WIDE backbone.

WIDE X6-Bone project





KR-JP Inter-operability test

Inter-operability check between *BSD and Linux implementation(July 15 2003 in 54th IETF)

- Linux(Korea)
 - Soongsil University
- NetBSD(Japan)
 - Fujitsu Laboratories, Ltd.
 - Fujitsu Limited
 - Nara Advanced Institute of Science and Technology
 - Nippon Telephone and Telegraph East Corporation
 - Matsushita Electric Industrial Co., Ltd.
 - Sony Computer Science Laboratories, Inc.
 - Information Services International-Dentsu, Ltd.
 - NoBUG: Northern Land BSD Users Group (Hokkai-do)
 - NBUG: Nagoya *BSD Users' Group



IETF Standardization

1999: 3 independent drafts were submitted

- Connectionless Multicast (Alcatel)
- Multiple Destination Option on IPv6 (Fujitsu)
- Small Group Multicast (IBM)

2000:

- 1st. BoF in 48th IETF
- Unified XCAST specification
 - Explicit Multicast Basic Specification
 - *draft-ooms-xcast-basic-spec-xx.txt*

2003:

- Preparing to start standard track discussion in RMT WG (Transport Area) in 56th IETF.

Conclusion

XCAST6 is new type of multicast

- Use list of unicast addresses as a destination of datagram.
- Suitable for private small group multicast
- Ultra scalable concerning with the number of multicast groups
- End-to-end deployment with semi-permeable capsule

2 inter-operable implementations for
Linux and *BSD

IETF standardization is just kicked off.

Links & Resources

XCAST incubation group

- <http://www.xcast-ig.org>

WIDE XCAST WG & X6-Bone

- <http://www.xcast.jp>

*BSD implementations

- <http://www.sourceforge.net/projects/xcast6>

Linux implementation

- <http://www.ipv6.or.kr/xcast/>